

CHALLENGER III

Multiple Processor Operations Manual

(Preliminary Version)

Introduction

This is a preliminary manual covering the fundamental operations of the Challenger III triple-processor CPU Board. The Model 510 CPU Board is a new concept in microcomputing allowing the user to utilize any of three processors. The operation of the triple-processor and understanding of it is definitely not trivial, but can be mastered if one has a good basic understanding of microcomputers. This manual covers fundamental routines allowing the user to switch back and forth between processors, and to run Z-80, 6800, and 6502 programs under the benefits of 6502 disk operating systems for program storage. It is a preliminary manual in the sense that we are planning to create a large library of utilities and programs making use of the 510 Board. This will greatly expand its usability, particularly for persons not versed in software or programming. We are very concerned about your interests in software for the 6800 and Z-80 and would like you to fill out the questionnaire at the back of this manual to guide us in developing and marketing software packages that can be of use to you.

TABLE OF CONTENTS

I	INTRODUCTION	
II	OVERVIEW OF THE 510	4-6
A	MANUAL SWITCHING	4-5
	6502 OPERATION	
	Z-80 OPERATION	
	6800 OPERATION	
B	SOFTWARE SWITCH OVERVIEW	5-6
	PIA DISCUSSION	
	SWAPPABLE RAM	
III	UTILITIES DISKETTE	7-24
A	COMMAND DESCRIPTION	7
	SYSTEM REQUIREMENTS	
B	SWITCHING ROUTINES	8
	GO ADDRESSES AND VECTORS	
C	TABLE OF SWITCH ROUTINES	9
D	EXECUTING AND HANDLING Z-80/8080 PROGRAMS	10-11
	LOADING Z-80/8080 SERIAL MONITOR	10
	STORING Z-80/8080 PROGRAMS ON DISK	
	LOADING Z-80/8080 PROGRAMS FROM DISK	
	SUMMARY OF Z-80/8080 OPERATION	11
E	Z-80/8080 MONITOR	12-16
	SYSTEM COMMANDS	12
	DESCRIPTION OF USEFUL SUBROUTINES	
	Z-80/8080 MONITOR LISTING	13-16
	Z-80/8080 MOVE ROUTINE LISTING	16
F	EXECUTING AND HANDLING 6800 PROGRAMS	17-19
	LOADING CHECKSUM LOADER AND DUMPER	17
	STORING 6800 PROGRAMS ON DISK	
	LOADING 6800 PROGRAMS FROM DISK	
	SUMMARY OF 6800 OPERATION	18
	USING THE 6800 UTILITIES	
	6800 MOVE ROUTINE LISTING	19
G	6800 MIK-BUG SIMULATOR	20-21
	SYSTEM COMMANDS	20
	IMPORTANT MEMORY LOCATIONS	21
H	A PROGRAMMER'S GUIDE TO THE SOFTWARE SWITCH	22-24
	ON THE PIA AND SWAPPABLE RAM	22-23
	SOFTWARE SWITCHING TABLE	23
	PROCESSOR SELECT SUBROUTINES	24
IV	EXAMPLES OF OPERATION	
V	68A2 MANUAL	
VI	6800 SPECIFICATIONS SHEET	
VII	QUESTIONNAIRE	

Challenger III CPU

The Model 510, or Challenger III, CPU Board features the 6502A, Z-80, and 6800 microprocessors. This CPU Board allows you to run programs for the 6502, Z-80, 6800, or 8080 on your Challenger III computer system. Operation of Z-80, 6800, and 8080 programs is fairly straightforward once you become completely familiar with the general operation of your computer system. Therefore it is very important that you become fluent in the operation of the 6502 portion of the system before you explore these other processors.

We will first discuss in general how the Challenger III CPU works before stating specific examples and presenting demonstration programs. The Challenger III is available in two configurations. One is the stripped-down or standard version with a processor-select switch. The other is a deluxe version, which has a software switch, a swappable RAM memory (i. e., RAM scratchpad), and a PIA port.

Let us first consider the operation of the minimal CPU, that is, the CPU with the selector switch. The basic Challenger III, or 510 CPU Board, is equipped with a single pole three position switch. The position of this switch specifies which processor is activated. The other processors are off and in the tristate condition. The 6502 position indicates that when a reset operation takes place by the reset switch, the processor will vector to the floppy disk bootstrap program, and the user can then elect to boot in the disk or go to 68A Monitor program. When the selector switch is in the Z-80 position, the restart occurs from location 0, so that when the user depresses the reset switch and releases it, the processor will start to execute code at location 00 (hex). The user must place his program at 00 (hex), or preferably place a JMP, or C3 (hex) command, to a location and page, in that order, to start up his program. Since users will most likely want to bring in their Z-80 and 6800 programs via the 6502 Disk Operating System, it is generally desirable to place their programs in high memory so that they can be co-resident with the 6502 DOS. For this reason, the user will generally want to put a JMP vector at location 00 for Z-80 programs. All standard 510 CPU Boards with or without software switching option are equipped with 68A Mod 4 PROM Monitor, such that when a user selects a 6800 and resets the computer, the processor will vector to the 68A4 Monitor program. The 68A Monitor provides the minimum requirements for manipulating and writing programs in the 6800, including the LOAD, PRINT, and GO commands.

The 68A documentation includes a discussion of how to convert MIKBUG-compatible programs to run under its system. MIKBUG is a Monitor ROM that has been marketed for a long time by Motorola. The MIKBUG Monitor was used in some of Motorola's early evaluation kits, and has been used in the SWTP 6800 computer.

Several independent software companies such as TSC have written extensive 6800 software packages utilizing the MIKBUG Monitor. The problem with MIKBUG compatibility is that MIKBUG was written for a small computer system. It is a 512-word Monitor residing at E000 to E200. It also makes use of a subroutine stack and scratchpad area at A000 up for 128 bytes. MIKBUG-dependent programs usually start at 0100 (hex) and have branches or subroutines in the MIKBUG Monitor at E000. They may also make use of the small amount of scratchpad memory at A000. The OSI system user has two alternatives when dealing with MIKBUG-dependent programs, that is, programs written to run the 6800 systems with the MIKBUG Monitor. He can follow the instructions provided with the 68A Monitor, which show how to change MIKBUG

subroutine calls from MIKBUG to the 68A. This basically means changing them from an address at E000 to one on page FF, and changing the A0 addresses to 1E addresses. This is a straightforward operation in fully documented programs such as those written by TSC, but is not so easy on undocumented or lightly documented programs. The other alternative is to utilize our MIKBUG simulator routines, which actually reside at E000 (hex). It will be necessary for the user to provide memory residing at these locations to operate these programs. It will minimally require that memory reside from E000 to E200 and A000 to A100, in addition to its normal memory. If he purchases a system with 44K of RAM or more, he will automatically have the memory at A000 (hex). In this case it will be necessary for him to add only a fully or partially populated 4K memory board at E000 up. With these boards and the MIKBUG simulator program in place, it will be possible to run MIKBUG-dependent 6800 programs without modification to the programs.

In place of the selector switch, the 510 Board with the software switch option utilizes a PIA output port to control certain parameters of the board. The processor select switch is documented on page 52 of the Comprehensive Information Package II. Two signal lines, or outputs, from the PIA at F700 are fed into a 74155 at IC-E10. This 74155 does a two-to-four-line decoding on output 0 to 3, which are fed into a 7475 latch. The three possible processor select codes are: 1-1, 6502; 1-0, 6800; 0-1, Z-80; and 0-0, all three processors tristated. Another PIA line, processor select, is used to actually change processors. This fires a series of one-shots causing all microprocessors to be reset, and then loads the new processor-select code into the 7475, bringing the processors out of reset, thus selecting the one and only one processor which is to be activated. The timing diagram is shown in the upper right corner of page 52 (Comprehensive Information Package II). An external reset switch performs a master reset, which always returns command to the 6502A when actuated. The master reset switch on the front panel always returns control to the 6502A processor when the software switch option is installed on the 510 CPU Board. To change processors, the program must first set up the new processor-select code on the PIA, which will then present this code to the input side of the 7475 latch. After this is done, the software must exercise the processor-select line on the PIA. This fires the one-shots, which reset the processors and load in the new processor-select code into the 7475, thus selecting the new processor and bringing it out of reset. Since both the 6502 and 6800 reset to their Monitor programs, the software change from a 6502 to a 6800 would normally bring a 6800 up in its Monitor program, which would then require additional manual intervention to run the user program. The Z-80, however, vectors to location 0, which can simply contain a JMP to the desired program. With this hardware alone, we can call up any specific Z-80 program from either the 6502 or the 6800.

To have the 6502 or 6800 call up and execute a specific program on the other processors, it is necessary to have additional hardware. This hardware is a swappable, or relocatable, RAM normally located at F200 up. It can be moved to FF00 up. This is accomplished by a single line on the PIA. When this line is actuated, this RAM becomes write-protected and is relocated in the address space of the computer to FF00, disabling the normal Monitor EPROM.

The obvious use for this RAM is to store new restart and interrupt vectors for either the 6502 or the 6800, such that when the

user performs a processor select, this RAM has been swapped up to the restart vector location and contains the restart vector to the program of interest. By invoking the swappable RAM, the user can (completely under program control) switch from a 6502 execution to that of a specific program under the 6800, and vice versa. A secondary purpose for this RAM is to allow the user to change his breakpoint or interrupt vectors at will, or from program to program, permitting practical utilization of interrupts on a 6502 or 6800 system.

Now that we have covered the operation of the Model 510 CPU Board in general terms, let us consider some specific utility programs and examples of switching from one processor to another.

CHALLENGER III UTILITIES DISKETTE

Included with this manual is a diskette containing utilities for handling the 6800 and Z-80 processors on the 510 CPU Board. The diskette contains a modified version of OS-65D in the sense that it does not boot in BASIC, the Assembler or Editor, but simply the Operating System, I/O Distributor and File Handlers. The operation is simply to put the diskette in the drive, reset the computer and type 0. It will come up with the Disk Operating System prompter A*. The operating system will make use of locations \$0000 through \$01FF and \$2180 through \$3180. Other memory locations will be undisturbed by the Disk Operating System.

The DOS has four additional commands. The VU automatically brings in the Challenger III Utilities Package (documented on the following pages). This occupies from \$3D70 TO \$3FFF. Thus it will run on a 16K Challenger III system. The Challenger III Utilities Package provides a necessary Z-80 Monitor, 6800 and Z-80 Memory Moves, and automatic GO routines between 6800, Z-80 and back for units with a software switch option. The automatic routines which switch between processors are not usable on systems without software switches.

Two commands which can be used only on units with a software switch and after Challenger III utilities have been brought in by the VU command are R6 and RZ. R6 returns to the 6800 operating under the 68A2 Monitor, and RZ switches the machine to the Z-80 operating under the Z-80 Serial Monitor. Remember of course that the standard OS-65D has an RM which returns to the 6502 65A PROM Monitor.

To summarize, we have a VU, which automatically boots in the Challenger III utilities package; RM, which returns command to the 65A Monitor under 6502; R6, which returns command to the 6800 under the 68A2 Monitor (only on units with the software switch, and after the VU command has been executed); and RZ, which returns command to the Z-80 system under the Z-80 Monitor (only on units with the software switch option and after the VU command has been executed).

Three additional tracks outside the operating system have useful utilities for the Challenger III system. On track 8 is another copy of the Challenger III Utilities Package, which is offered for user modification if desired. Track 9 contains a MIKBUG simulator program for use on the 6800 portion of the system. It is identical in operation to the very popular Motorola 6800 MIKBUG Monitor Program, except it uses our standard ACIA I/O instead of the PIA I/O. With this Monitor up and running, the Challenger system will be functionally equivalent to a 6800 computer which runs under MIKBUG.

Requirements to run the MIKBUG simulator are a memory board with at least 512 bytes of memory at location E000 (recommended: a partially or fully populated 420 board addressed at E000) and at least 128 bytes at A000. A 48K computer will automatically have this; if not, another 420 Board will have to be strapped here.

Track 10 has a 6800 load, dump, and print data string utilities package. When this is used in conjunction with the 65A2 Monitor, it provides load and dump capabilities in the popular Motorola S-1 checksum format, and a P-Data subroutine similar to the one found in MIKBUG.

The various modules of the Challenger III utilities diskette are covered in detail in the following sections, and at the end of the manual with specific examples of use.

510 Utilities Package

The 510 Utilities Package is a collection of routines which allow convenient operation in switching between 6502, 6800, and Z-80. This utilities resides from \$3D70 to \$3FFF. It is booted via the (VU) with the modified OS-65D operating system on your Z-80 utilities disk and is also provided as a separate track for use and modification under other circumstances.

The utilities package has nine modules. The first six modules are Processor Select routines for software switch units only. By simply jumping to the go address specified, the user will switch from the specified processor to the specified processor. For example, entering the routine at 3DAA will switch the machine from the 6800 to the 6502. If a routine is called improperly, that is, if the machine is currently running as a Z-80 and the user calls the routine at 3D70, it will crash. So it is important to be careful in selecting these routines.

The result of these six routines is that the new processor comes up in its Monitor program. In the case of coming up as a Z-80, the Monitor program is located at 3E00 up, as documented separately. In conjunction with the Z-80 Monitor, the routines which select the Z-80 place a JMP 3E00 at location 00, which remains. Following the Z-80 Monitor is the Z-80 Memory Move routine, which can be executed only by the Z-80 processor, of course. With this Move routine, under the Z-80 Monitor, for instance, the user will place the low address and high address of the start of the move (\$3F4F and \$3F50), the low address and high address of the end of the block to be moved (\$3F51 and \$3F52), and the starting destination address low, then high (\$3F53 and \$3F54). He then jumps to the Move program at \$3F55, which performs the Move operation, and returns command to the Z-80 Monitor. The user must be careful not to move the program in on top of the Z-80 Monitor, or the Move program, or a crash will occur.

Likewise there is a Memory Move for the 6800 located at \$3F79 up. This Memory Move can be operated on by the 68A2 Monitor, for instance. Remember that the 6800, unlike the 6502 and Z-80, uses a syntax of high address and low address. So at location \$3F79 put the high address of the beginning of the move, and at \$3F7A put the low address. Likewise at \$3F7B and \$3F7C put the end of the block to be moved, and at \$3F7D and \$3F7E put the beginning destination address in the order high address-low address. Be very careful to remember the address difference between the 6800, Z-80, and 6502. Confusion here leads easily to computer crashes.

510 UTILITIES

Go Address	Processor (Old)	Processor (New)	
3D70	6502	6800	
3D88	6502	Z-80	
3DAA	6800	6502	
3DBF	6800	Z-80	
3DD4	Z-80	6502	
3DEA	Z-80	6800	
<hr/>			
3E00	Z-80	Monitor	
<hr/>			
3F55	Z-80	Mover	3F4F/50 Source Address 3F51/52 End Address 3F53/54 Dest. Addr.
<hr/>			
3F7F	6800	Mover	3F79/7A Source Address 3F7B/7C End Address 3F7D/7E Dest. Addr.

Executing and Handling Z-80 and 8080 Programs

Call in the the 510 utility via the VU command. On units with the software switch, then type RZ. On units with the mechanical switch, return to the 65A via an RM, then place the JMP 3E00 via L0000 C3 00 3E R. Then hold in the reset switch and rotate the processor switch to Z-80; release the reset switch. You should now see the Z-80 Monitor prompter "?". The user can then enter programs and run them via the Monitor. If large programs available on paper tape or cassette are to be entered, it would be desirable to write a loader program which could then be used to load other programs. Another approach is to write such loaders for the 6502, thus making use of the 6502 Assembler.

Storing Z-80 Programs on Disk

The 6502 Disk Operating System uses all of Page 0 and 1 (of necessity) and \$2200 to \$3200. 6502 modules such as BASIC or the Assembler use additional memory. To conserve memory a DOS-only operating system is present on the Challenger III utilities diskette. Programs which do not use \$0-\$200 and \$2200 to \$3200 can be directly stored and retrieved from this disk by using the C (=CALL) and S (=SAVE) commands. This is accomplished on units with the mechanical switch by holding in reset while rotating back to the 6502. Then reboot the 6502 DOS (just to be safe) and save the program of interest on disk. If the program of interest uses memory locations which are used by the 6502 DOS, use the move utility at the end of the Z-80 Monitor to move the program to high memory, for example, at \$4000 up then save it via the 6502 DOS. Likewise the move routine can be used to place a program back in position after being brought in via the 6502 DOS.

The procedure for getting back to the 6502 from the Z-80 on units with the software switch is either to enter the routine at \$3DD4 or preferably just reset the machine.

Calling Z-80 Programs from Disk

Under the 6502 DOS call in the utilities package and your program. On software switch units type RZ. On mechanical switch units exit the DOS via an RM. Place a JMP to the Z-80 Monitor at location 00 via L0000 C3 00 3E R, then hold the reset switch in, switch to Z-80 and release. In either case, once you are in the Z-80 Monitor you can go to the program via the Monitor's GO command. Entering the Z-80 Monitor can be avoided by setting the JMP at location 00 directly to the start of the program. However it is generally desirable to have the Z-80 restart set for the Z-80 Monitor so that the Z-80 can be restarted in case of a crash.

Summary of Z-80 Operation

To enter programs:

1. Bring in 6502 DOS
2. Call in Z-80 Monitor
3. Call in Challenger III utilities
4. Switch to Z-80 (on manual switch units first set up C3003E)
5. Move program in place if necessary; if new program, key in via Monitor

To save programs:

1. Move program to high memory if necessary (above \$4000)
2. Switch back to 6502
3. Reboot 6502 DOS and save

Z-80 Monitor

The Z-80 Monitor is a short machine code program that is provided to reside at 3E00 up for approximately one page. This program offers elementary memory examine, modify and program execute commands on a serial system using a standard ACIA port as the I/O device.

To operate this program one must first place it in memory typically by use of the 6502 Disk Operating System. The user must place a JMP 3E00 at location 00 in memory (C3 00 3E). Then the user must either manually reset the processor as a Z-80 or go through the software switch and convert the system to a Z-80. Keep in mind that the 6502 DOS and BASIC and other large programs will wipe out the JMP at location 00 if executed, so that the JMP placed at location 00 up should be done after exiting the DOS or BASIC, just before the user actually turns on the Z-80. The Z-80 Monitor comes up with a ? prompt. The user types in a four-digit hexadecimal address, for example, AAAA. The monitor then displays the contents of that address as two hex digits. The user can then optionally modify those locations by typing a two-digit hex number. As soon as the two-digit hex number is typed in, the monitor will automatically display the next address and its contents. The user can examine memory by simply typing <return> immediately after the contents of the memory location are displayed. By successively typing <return> he will display consecutive memory locations. When the user types a ^, the Monitor will show the memory location previous to the one shown.

The user can exit the display mode at a particular address by typing an R. This returns the Monitor to the command mode, where it expects to see an address again. Thus the user can look at a few locations, type an R, type in a new address, and examine locations there. To execute a program the user must type in the starting address of the program, then an RG, which will cause program execution at the specified address.

The Z-80 Monitor contains an INPUT subroutine with echo located at 3E91. This subroutine puts ASCII character in the Accumulator upon exit The Accumulator is modified. Another useful routine in the Z-80 Monitor is the Character Output Routine. This routine is located at 3E90 and returns the ASCII value, outputting the ASCII equivalent of the Accumulator. Other useful routines are as follows :

3F10 LEGAL - prints a ? if (A)<>ASCII for 0-->9 or A-->F
also sets carry
if (A)=ASCII for 0-->9 or A-->F the sub returns the hex
value with the carry=0

3F20 PROMPT- prints a ? and sets the carry flag (A lost)

3F44 CR-LF - prints carriage return-linefeed (A lost)

3F38 DIGIT - converts hex 0-->9 or A-->F to ASCII code (uses A only)

3EAB DISP DATA- prints </nn>, where nn= contents of
memory location pointed to by \$FE(low) \$FF(high)

Stack Pointer initialized at \$100.

SERIAL Z-80/8080 MONITOR

```

3E00 318001  START  LXI SP, $180
3E03 210000      LXI H, $0
3E06 22FE00      SHLD $FE
3E09 2100FC      LXI H, $FC00
3E0C 3E03        MVI A, 3
3E0E 77          MOV M, A
3E0F 3EB1        MVI A, $B1
3E11 77          MOV M, A
3E12 CD443F      CALL CR-LF
3E15 CD2C3F      CALL PROMPT

3E18 CD443F  AD      CALL CR-LF
3E1B 21FF00      LXI H, $FF
3E1E 0602        MVI B, 2
3E20 0E02  REP2   MVI C, 2
3E22 CD913E  REP1   CALL INCH
3E25 FE47      CPI G
3E27 CA8D3E      JZ GO
3E2A FE52      CPI R
3E2C CA183E      JZ AD
3E2F CD113F      CALL LEGAL
3E32 DA183E      JC AD
3E35 0D          DCR C
3E36 CA443E      JZ SKIP1
3E39 17          RAL
3E3A 17          RAL
3E3B 17          RAL
3E3C 17          RAL
3E3D 77          MOV M, A
3E3E 3E3F        MVI A, 0
3E40 B9          CMP C
3E41 C2223E      JNZ REP1
3E44 B6  SKIP1   ORA M
3E45 77          MOV M, A
3E46 2B          DCR B
3E48 C2243E      JNZ REP2
3E48 CDAB3E      CALL DISP DA

3E4E 0E02  DA     MVI C, 2
3E50 CD913E  REP3   CALL INCH
3E53 FE52      CPI R
3E55 CA183E      JZ AD
3E58 FE5E      CPI $5E
3E5A CA833E      JZ BACK
3E5D FE0D      CPI $D
3E5F CA793E      JZ NEXT
3E62 CD113F      CALL LEGAL
3E65 DA183E      JC AD
3E68 0D          DCR C
3E69 CA773E      JZ SKIP2
3E6C 17          RAL
3E6D 17          RAL
3E6E 17          RAL
3E6F 17          RAL
3E70 47          MOV B, A
3E71 3E00        MVI A, 0
3E73 B9          CMP C
3E74 C2503E      JNZ REP3

```

3E77	B0	SKIP2	ORA B
3E78	77		MOV M, A
3E79	23	NEXT	INX H
3E7A	CDCC3E		CALL DISP AD
3E7D	CDAB3E		CALL DISP DA
3E80	C34E3E		JMP DA
3E83	2B	BACK	DCX H
3E84	CDCC3E		CALL DISP AD
3E87	CDAB3E		CALL DISP DA
3E8A	C34E3E		JMP DA
3E8D	2AFE00	GO	LHLD \$FE
3E90	E9		PCHL
3E91	3A00FC	INCH	LDA \$FC00
3E94	1F		RAR
3E95	D2913E		JNC INCH
3E98	3A01FC		LDA \$FC01
3E9B	E67F		ANI \$7F
3E9D	F5	OUTCH	PUSH PSW
3E9E	3A00FC	WAIT	LDA \$FC00
3EA1	1F		RAR
3EA2	1F		RAR
3EA3	D29E3E		JNC WAIT
3EA6	F1		POP PSW
3EA7	3201FC		STA \$FC01
3EAA	C9		RET
3EAB	3E2F	DISP DA	MVI A, '/'
3EAD	CD9D3E		CALL OUTCH
3EB0	2AFE00		LHLD \$FE
3EB3	7E		MOV A, M
3EB4	F5		PUSH PSW
3EB5	1F		RAR
3EB6	1F		RAR
3EB7	1F		RAR
3EB8	1F		RAR
3EB9	CD383F		CALL DIGIT
3EBC	CD9D3E		CALL OUTCH
3EBF	F1		POP PSW
3EC0	CD383F		CALL DIGIT
3EC3	CD4D3E		CALL OUTCH
3EC6	3E20		MVI A, \$20
3EC8	CD9D3E		CALL OUTCH
3ECB	C9		RET
3ECC	CD443F	DISP AD	CALL CR-LF
3ECF	22FE00		SHLD \$FE
3ED2	0602		MVI B, 2
3ED4	0E02		MVI C, 2
3ED6	7C	REP4	MOV A, H
3ED7	05		DCR B
3ED8	CADF3E		JZ SKIP3
3EDB	1F		RAR
3EDC	1F		RAR
3EDD	1F		RAR
3EDE	1F		RAR
3EDF	CD383F	SKIP3	CALL DIGIT
3EE2	CD4D3E		CALL OUTCH
3EE5	0D		DCR C

3EE6	C2D63E		JNZ	REP4
3EE9	0602		MVI	B, 2
3EEB	0E02		MVI	C, 2
3EED	7D	REP5	MOV	A, L
3EEE	05		DCR	B
3EEF	CA063F		JZ	SKIP4
3EF0	C3023F		JMP	EXTRA
3EF3	00		NOP	
3EF4	00		NOP	
3EF5	00		NOP	
3EF6	00		NOP	
3EF7	00		NOP	
3EF8	00		NOP	
3EF9	00		NOP	
3EFA	00		NOP	
3EFB	00		NOP	
3EFC	00		NOP	
3EFD	00		NOP	
3EFE	00		NOP	
3EFF	00		NOP	
3F01	00		NOP	
3F02	1F	EXTRA	RAR	
3F03	1F		RAR	
3F04	1F		RAR	
3F05	1F		RAR	
3F06	CD303F	SKIP4	CALL	DIGIT
3F09	CD9D3E		CALL	OUTCH
3F0C	0D		DCR	C
3F0D	C2ED3E		JNZ	REP5
3F10	C9		RET	
3F11	FE30	LEGAL	CPI	\$30
3F13	DA2C3F		JC	FAIL
3F16	FE3A		CPI	\$3A
3F18	DA273F		JC	OK
3F1B	FE41		CPI	\$41
3F1D	DA2C3F		JC	FAIL
3F20	FE47		CPI	\$47
3F22	D22C3F		JNC	FAIL
3F25	D607		SUI	7
3F27	E60F	OK	ANI	\$F
3E29	37		STC	
3E2A	3F		CMC	
3F2B	C9		RET	
3F2C	3E20	FAIL	MVI	A, \$20 (PROMPT)
3F2E	CD9D3E		CALL	OUTCH
3F31	3E3F		MVI	A, '?'
3F33	CD9D3E		CALL	OUTCH
3F36	37		STC	
3F37	C9		RET	
3F38	E60F	DIGIT	ANI	\$F
3F3A	F634		ORI	\$30
3F3C	FE3A		CPI	\$3A
3F3E	DA433F		JC	EXIT
3F41	C607		ADI	7
3F43	C9	EXIT	RET	
3F44	BE0D	CR-LF	MVI	A, \$D
3F46	CD9D3E		CALL	OUTCH

3F49 3E0A
3F4B CD9D3E
3F4E C9

MVI A, 0A
CALL OUTCH
RET

Z-80 Mover

Source address low @ 3F4F, high @ 3F50
End address low @ 3F51, high @ 3F52
Destination address low @ 3F53, high @ 3F54

```
3F55 2A513F      LHL      ; END ADDRESS-->(HL)
3F58 EB          XCGH      ; END ADDRESS-->(DE)
3F59 2A4F3F LOOP LHL      ; SOURCE ADDRESS-->(HL)
3F5C 7E          MOV A,M   ; (M)-->(A)
3F5D 2A533F      LHL      ; DESTINATION ADDR. -->(HL)
3F60 77          MOV M,A   ; (A)-->(M)
3F61 23          JNX H    ; (HL+1)-->(HL) (=INCREMENTS...
3F62 22533F      SHLD     ; DESTINATION ADDRESS & STORES IT)
3F65 2A4F3F      LHL      ; SOURCE ADDRESS-->(HL)
3F68 7A          MOV A,D   ; END ADDRESS HIGH-->(A)
3F69 BC          CMP H    ; ARE WE ON UPPER LIMIT PAGE?...
3F6A C2723F      JNZ SKIP ; ... IF NOT, SKIP TO INCR. SOURCE ADDR.

3F6D 7B          MOV A,E   ; END ADDR. LOW-->(A)
3F6E BD          CMP L    ; ARE WE DONE?
3F6F CA123E      JZ EXIT  ; YES-RETURN TO MONITOR

3F72 23          SKIP INX H   ; (INCREMENTS SOURCE ADDRESS...
3F73 224F3F      SHLD     ; ... & STORES IT)
3F76 C3593F      JMP LOOP ; NOT FINISHED LOOP BACK
```

Executing and Handling 6800 Programs

The 6800 portion of the 510 Board has its own 256-byte PROM Monitor located at FF00. To turn on the 6800 and enter the 68A2 Monitor on mechanical switch units, simply hold the reset switch in, rotate the switch to the 6800 position and release the reset switch. The terminal should output a carriage return-linefeed and be in the 6800 Monitor. On units with the software switch call in the utilities package via the DOS and type RG, which will bring up the 6800.

Once in the 68A2 Monitor, programs can be entered manually if desired. If long programs are available on cassette or paper tape, it would be desirable to create a loader program to bring these in. A Motorola "S" format checksum tape loader and dumper is included in the 6800 utilities package. This format is almost universally used for all paper tape and most audio cassette programs for the 6800.

Storing 6800 Programs on Disk

The 6502 Disk Operating System uses all of Page 0 and 1 (of necessity) and \$2200 to \$3200. 6502 modules such as BASIC or the Assembler use additional memory. To conserve memory a DOS-only operating system is present on the Challenger III utilities diskette. Programs which do not use \$0-\$200 and \$2200 to \$3200 can be directly stored and retrieved from this disk by using the C (=CALL) and S (=SAVE) commands. This is accomplished on units with the mechanical switch by holding in reset while rotating back to the 6502. Then reboot the 6502 DOS (just to be safe) and save the program of interest on disk. If the program of interest uses memory locations which are used by the 6502 DOS, use the 6800 move utility to move the program to high memory, for example, at \$4000 up, then save it via the 6502 DOS. Likewise the move routine can be used to place a program back in position after being brought in via the 6502 DOS.

The procedure for getting back to the 6502 from the 6800 on units with the software switch is either to enter the routine at \$3DAA or preferably just reset the machine.

Calling 6800 Programs from Disk

Under the 6502 DOS call in the utilities package and your program. On software switch units type RG. On mechanical switch units then hold the reset switch in, switch to 6800 and release. In either case, once you are in the 68A Monitor you can go to the program via the Monitor's GO command. Entering the 68A Monitor can be avoided by setting up the swappable RAM (as covered in a later section) directly to the start of the program. However it is generally desirable to have the 6800 restart set for the 6800 Monitor so that the 6800 can be restarted in case of a crash.

Summary of 6800 Operation

To enter programs:

1. Enter 6502 DOS.
2. Call in Challenger III Utilities.
3. Call in 6800 Program from Disk.
4. Switch to 6800.
5. Move Program in place if necessary; if new program, key in via 6802 Monitor.

To save programs:

1. Move Program to high memory if necessary (above \$4000).
2. Switch back to the 6502.
3. Reboot 6502 DOS to save.

6800 Utilities

Track 10 of the Challenger III diskette contains a 6800 load-dump and p-data routine that is documented in the 68A manual. To utilize this program, simply type C1D00=10,1, utilizing the instructions in the 68A writeup. The 6800 Memory Mover is part of the Challenger III utilities. Its operation is documented there. Finally the 6800 MIKBUG Simulator is provided on track 9, and its use is discussed in the following section.

6800 Move File (relocatable)

```

1E80  3F7F  FE  LDX
1E81  3F80  1E
1E82  3F81  02
1E83  3F82  A6  LDA, X
1E84  3F83  00
1E85  3F84  FE  LDX
1E86  3F85  1E
1E87  3F86  06
1E88  3F87  A7  STA, X
1E89  3F88  00
1E8A  3F89  00  INX
1E8B  3F8A  FF  STX
1E8C  3F8B  1E
1E8D  3F8C  06
1E8E  3F8D  FE  LDX
1E8F  3F8E  1E
1E90  3F8F  02
1E91  3F90  BC  CPX
1E92  3F91  1E
1E93  3F92  04
1E94  3F93  27  BEQ
1E95  3F94  06
1E96  3F95  00  INX
1E97  3F96  FF  STX
1E98  3F97  1E
1E99  3F98  02
1E9A  3F99  20  BRA
1E9B  3F9A  E7
1E9C  3F9B  7E  JMP
1E9D  3F9C  FF  (mod for 510 system)
1E9E  3F9D  B2
1E9F  ---

```

	Source	High Limit	Destination
High	3F79	3F7B	3F7D
Low	3F7A	3F7C	3F7E

6800 MIKBUG Simulator

The MIKBUG Simulator is on track 9 of the Challenger III utilities diskette. It is a 512-byte program residing from \$E000 to E1FF. The MIKBUG simulator is functionally equivalent to the popular Motorola MIKBUG Monitor, with the exception that the Motorola Monitor uses a PIA interface in conjunction with some discrete logic to perform a serial I/O, while in the MIKBUG Simulator the standard 6850 ACIA port at FC00 is utilized. Both MIKBUG and the MIKBUG Simulator reside at E000 and require a subroutine stack of at least 128 bytes at A000 up. The MIKBUG Simulator program is useful because most hobbyist-level software written for the 6800 has been written on MIKBUG-based 6800 machines, and generally this software makes heavy use of subroutines located within the MIKBUG Monitor, and also makes use of the 128-byte scratchpad RAM normally located at \$A000. Therefore having the capability of configuring your Challenger III system to be virtually identical to a MIKBUG-based 6800 system may be of value.

To accomplish complete compatibility and run MIKBUG, it is necessary to have a total of 512 bytes of RAM memory at \$E000 up. The recommended way of achieving this is to partially populate or fully populate a Model 420C Memory Board and address it for \$E000 up and place it in the system. It will also be necessary to have at least 128 bytes of RAM memory at \$A000. If the machine is a 48K computer, it will automatically have memory here. If it is less than 44K, it will be necessary to add a fully or partially populated 420C Memory Board at this location also. So if memory is at \$A000 for at least 128 bytes, and at \$E000 for at least 512 bytes, the user can call in the MIKBUG Simulator on track 9 by typing CE000=09,1. Then either manually or under software control go to the 68A2 Monitor. The MIKBUG Simulator Monitor is entered at \$E0D0. So from 68A2 type L1F2E E0D0 RG. The machine will then output a *, which is the MIKBUG prompter.

The MIKBUG Simulator is completely identical in operation to the standard MIKBUG. The operation of MIKBUG is well documented in Motorola Engineering Note 100. The MIKBUG Monitor has five basic commands: L (LOAD), M (MEMORY), P (PUNCH), G (GO), and R (REGISTER DUMP). Typing an L with the prompter on the screen causes the MIKBUG Simulator to go into the checksum loader routine. This routine will accept the popular Motorola S-1 checksum format via the ACIA port. However, if one is not using a Teletype as an input device, it will be necessary to modify this routine for high-speed paper tape or audio cassette. It would be advisable to utilize the 6800 load-save routine on track 10 instead, or S-1 format checksum loads. The load is terminated by a CHECKSUM error, which will cause a ? and return to the command mode or an S9, which terminates the record.

The most useful command of the MIKBUG is the Memory Examine and Modify command. To utilize this type M followed by a four-digit address in hex. The MIKBUG then reprints the address followed by the contents of that location. To simply examine locations type <return>s. To modify a location type <space-bar> followed by two characters to be loaded into that location. To exit the Memory Load and Examine routine, type <space-bar> <return>.

MIKBUG also contains a Punch command which lists out a block of memory in the Motorola S-1 format. Before the Punch command is executed, it is necessary to use the Memory command to load the starting and ending addresses of the Punch. The starting address must be loaded into address A002 and A003 (high-low), and the ending

address into address A004 and A005 (high-low). Once these four locations have been loaded, type P and the MIKBUG Monitor will printout a checksum record of the memory.

Another command of the MIKBUG Monitor is the R command (REGISTER DISPLAY). Type an R and the Monitor puts out the current contents of the Condition Code Register, the Accumulators (A and B), Index Register (high and low), and the Go stack or interrupt stack. This Go stack is just like the stack in the OSI 68A2 Monitor, that is, these register values will be placed into the microprocessor upon execution of a Go command and will be reloaded into the monitor upon execution of a software interrupt, so that it will perform both go and breakpoint functions. The register values are stored in memory locations \$A043 (Condition Code Register), \$A044 (Accumulator A), \$A045 (Accumulator B), \$A046/47 (Index Register), \$A048/49 (Program Counter), and \$A04A/4B (Stack Pointer), so that the user is free to modify the status of all registers, including most importantly the program counter, before starting a program after breakpoint, or initially, by way of the Go command.

The final instruction is a Go command. By typing a G, the contents of A043 to A04B are loaded into the corresponding microprocessor registers, and program execution is then resumed. It is always necessary for the user to specify at least Program Counter at \$A048/49 and a Stack Pointer at \$A04A/4B before typing a G to execute a program.

We do not recommend that the user use MIKBUG under normal circumstances as a Monitor. We believe that the 68A2 Monitor is just as convenient as and certainly more accessible to the user. Also, the MIKBUG Monitor's breakpoint routine is useless unless its breakpoint vector has been loaded into the swappable RAM memory, and this has been swapped up to high memory. Memory should also be swapped to set the MIKBUG Simulator entry point E0D0, when it is used with programs making use of the MIKBUG reset and interrupt vectors. MIKBUG should be used in situations where MIKBUG must be resident in memory or the program in order to operate, and the user does not wish to modify the program to relieve its dependence on MIKBUG, and in situations where the swapping RAM must be swapped up to high memory, because under those circumstances the OSI 68A2 Monitor will not be usable because it is deactivated when the swappable RAM is in high memory (from FF00 up).

A Programmer's Guide to the Software Switch

As stated earlier, the software processor switch is controlled by the 6820 PIA at F700 up. On Challenger III CPUs equipped with the software switch, this PIA controls the selection of microprocessors and also optionally allows the user to exchange a 128-byte RAM normally located at F200 for the Monitor PROM located at FF00 up. Thus the user can pre-load this scratchpad RAM at F200 with new restart and/or interrupt vectors, then swap it up to replace the Monitor PROM. Hence it is called "swappable RAM."

Operation of the software switch to select processors and swap the RAM is possible manually by simply keying in the appropriate codes to the PIA control registers. The following table shows how this is accomplished. First the user enters \$04 into address F701. He then enters the appropriate processor select code into location F700. He loads a 00 into location F701 and an FF into F700. The action of loading the FF into F700 will trigger the Processor Selector switch into operation, which will then select the new processor and/or swap up the RAM at F200. The Processor Select codes are shown directly below, with and without swapping the RAM. This table means that if we put a 9F in for the PS code, this routine will force the machine to switch over to the 6800, no matter what the current state of the machine in. Likewise, if we load an 8F, the machine will revert to a 6800, and we use locations in the 128-word scratchpad memory for restart and interrupt vectors, regardless of the current state of the machine. So by this simple procedure, the user can manually switch between processors and swap up the RAM memory. The execution of this routine while utilizing the Processor Select code for the same procedure should have no effect other than resetting the processor. Using this routine with the Processor Select code for the processor you are currently using and the RAM swap will cause a reset with the RAM being switched up to high memory. This is caused when bit 4 of the Processor Select code is low. In that case the RAM is moved up to FF00 up, and the PROM Monitors for the 6800 at FF00 up and the 6502 disk bootstrap PROM at FF00 up are disabled. The RAM is also write protected when it is in this condition, so it is not possible for the user to change locations within the RAM when the RAM is at FF00 up. It is still possible to use the 65A Monitor after the RAM has been swapped, but it is not possible to use the 68A Monitor after the RAM is swapped, because this swap disables the Monitor.

When the RAM memory has been swapped up to FF00 up by having bit 4 of the processor select code high, the vectors contained at locations F2FA to F2FF are utilized. This means that in order to put in a new restart vector for the 6502, the user first writes the restart vector into location F2FC, and then swaps this RAM up to high memory. Remember that all 6502 vectors are in the form "low-order address, page," while all 6800 vectors are in the form "page, address."

The following page lists machine executable subroutines for switching from one processor to any other processor with or without RAM swap. These are three separate routines, one each for the 6502, 6800, and Z-80. In each case, the code is completely relocatable. For instance, for the 6800 program to return control to the 6502, the user simply places the 6800 processor select subroutine at the end of his program as listed, and substitutes a DF for the processor select code. When this is executed by the 6800, it will return command to the 6502. To perform the same function, but go to a specific address

other than the Monitor of the 6502, the 6800 programmer would first have to load the desired restart address into location F2FC low and F2FD high of the swappable RAM, with this RAM FC00, and then execute the 6800 Processor Select subroutine with a CF in the Processor Select code instead of a DF.

Software Switch

Manually through Monitor

```
04--->F701
PS--->F700
00--->F701
FF--->F700
```

PS	NO RAM SWAP	WITH RAM SWAP
Z-80	5F	4F
6800	9F	8F
6502	DF	CF

Vectors	6502	6800
	65A	68A2

PROM		
Restart	FE35	FFE7
NMI	0130	1FE0
IRQ	01C0	1FD0
Software	----	FFA8

RAM		
Restart	<F2FC>	<F2FE>
NMI	<F2FA>	<F2FC>
IRQ	<F2FE>	<F2F0>
Software	----	<F2FA>

PROCESSOR SELECT SUBROUTINES

For 6502--

```

A2 00      ENTRY  LDX 0
0E 01 F7   STX$F701  SELECT DDRA
CA         DEX
0E 00 F7   STX$F700  PORT A=OUTPUT
A0 04     LDY 4
0C 01 F7   STY$F701  SELECT DATA REG
0E 00 F7   STX$F700  SET ALL DATA HIGH
*A0 XX     LDY PS    PROCESSOR CODE
0C 00 F7   STY$F700
D0 FE     TRAP   BNE TRAP  WAIT FOR NEW PROCESSOR
    
```

For 6800--

```

CE F7 00   ENTRY  LDX $F700  PIA BASE ADDRESS
06 00     LDA A 0
A7 01     STA A 1, X  SELECT DDRA
4A         DEC A
A7 00     STA A 0, X  PORT A=OUTPUT
C6 04     LDA B 4
E7 01     STA B 1, X  SELECT DATA REG
A7 00     STA A 0, X  SET ALL DATA HIGH
*C6 XX     LDA B PS    PROCESSOR CODE
E7 00     STA B 0, X
3E         WAI      WAIT FOR NEW PROCESSOR
    
```

For Z-80--

```

21 00 F7   LXI H $F700
EB         XCHG
21 01 F7   LXI H $F701
3E 00     MVI A 0
77         MOV M, A    SELECT DDRA
3D         DCR A
12         STAX D     PORT A=OUTPUT
47         MOV B, A
3E 04     MVI A 4
77         MOV M, A    SELECT DATA REG
70         MOV A, B
12         STAX D     SET ALL DATA HIGH
*3E XX     MVI A PS    PROCESSOR CODE
12         STAX D
76         HLT      WAIT FOR NEW PROCESSOR
    
```

PS CODE--

PROCESSOR	NO RAM SWAP	WITH RAM SWAP
Z-80	5F	4F
6800	9F	8F
6502	DF	CF

*Shown in Immediate Mode, but may use any non-indexed addressing mode without disrupting program flow.

CHALLENGER III MULTIPLE PROCESSOR OPERATION EXAMPLES

```
D/M?D
A*VU
A*k
?
0000/C3
0001/00
0002/3E R

3DD4/21 R
D/M?D
A*k
P 0000
C3 00 3E 4C 0E 0A 0E 0F
04 11 00 18 20 30 20 20

L 1F2E 3DAA
LVMD
A*
```

Example 1. Challenger III utilities diskette was put in the disk drive, the machine was reset, the user typed a D. The operating system was booted in, returning a prompter A*. The user typed a VU, bringing in the Challenger III utilities. He then typed an RZ. The Z does not print, because the machine immediately switched to the Z-80, showing the Z-80 prompter "?". The user then typed in a four-character address (location 0), and examined that location. He typed <return> again, which outputted the following location. He typed <return> again, which outputted location 2, typed R to exit the memory examine routine. He then typed in address 3DD4, then RG. The G did not print, because the user went to address 3DD4, which immediately switched the machine back to the 6502. We then see the 6502 Monitor prompter D/M?. The Disk Operating System was booted in again. This time the user typed R6, which immediately switched the machine over to the 6800, the prompter being cr-1f. The user then printed out memory location from 0 on, he then loaded 1F2E (the go locations), with the address of the routine to bring him back to the 6502 (i. e., 3DAA). He then typed an RG, which does not print. The G brought the unit back to the 6502 DOS. This routine reinitializes the ACIA, so that the first character output may not be correct. Hence, instead of getting D/M? we got VM. He then typed D, which rebooted the operating system, and brought him back to the beginning. So in this example the user went from the 6502 to the Z-80, back to the 6502, then to the 6800, and back to the 6502.

Example 2. The user is in the DOS and calls in track 10, which is the load and save of the 6800. He then typed an R6, which took him to the 6800. The under the 68A2 Monitor, he loaded the start and end address of an S-1 format dump. He then loaded the GO vectors at 1F2E for the S-1 format dump program, and set the stack at 1F80. He typed RG (R does not print). The program at 1D70 then outputed an S-1 format checksum dump and exited back to the 68A2 Monitor

```

CID00=10,1
A*R
L 1E02 0000 0020
L 1F2E 1D70 1F80
G
SI 130000C3003E4C0E0A0E0F0411001820302020AD
SI 13001001300100202020202034323420202020F0
SI 04002000DB

```

Example 3. The user is in the 6502 DOS and calls into the MIKBUG Simulator on track 9. He then typed R6, which converted the machine to the 6800 under the 68A2 Monitor. He then loaded vector E0D0 into the 68A2 GO address at 1F2E and typed RG, which put him in the MIKBUG simulator and outputed the * prompter. He then executed a memory command in the MIKBUG Simulator, followed by address 0000. The monitor reported the contents of that location, and the user typed any key but a <space-bar> to cause the monitor to go through consecutive memory locations. Finally after address 3, the user typed a <space-bar> <return> and then used the REGISTER command, which outputed the current contents of the GO stack. The user then used a MEMORY command to modify location A045, A046, and A047, thus changing command then shows these changes. Finally the user changed the Program Counter contents to point back to the 68A2 monitor, and types a G (which does not print), returning the 6900 to the 68A2 monitor. He then prints out location 0, loads it with the sequence 01 through 0A, then examines location 200 and observes the data there. He then loads location 3F79 up with the start of move, end of move, and start of destination for the 6800 move portion of the Challenger III utilities. He then sets up the GO vector and subroutines stack at 1F2E and types a G. The move is nearly instantaneous and when it returns command to the 68A2 monitor, the user examines location 200 and verifies that the move did in fact occur.

```

A*CE000=09,1
A*R
L 1F2E E0D0

*M 0000
*0000 C3 .
*0001 00 .
*0002 3E .
*0003 4C
*R 24 24 24 2424 2424 A042
*M A043
*A045 24 10
*A046 24 20
*A047 24 30
*A048 24
*R 24 24 10 2030 2424 A042
*M A048
*A048 2C FF
*A049 2C A8
*A04A 2C
*A04B 2C
*
P 0000
80 80 80 80 80 80 80 80

L 0000 01 02 03 04 05 06 07 08
P 0200
80 80 80 80 80 80 80 80

L 3F79 0000 0010 0200
L 1F2E 3F7F 1F80
G
P 0200
01 02 03 04 05 06 07 08

```

```

D/M?D
A*VU
A*R
?
0100/90 10
0101/90 20
0102/90 30
0103/62 40
0104/FF 50
0105/FO 60
0106/FF 70
0107/90 80
0108/90 90

0109/91 ?
0300/90
0301/90
0302/90
0303/90

0304/90
0305/90
0306/90 ?
3F4F/FF 00
3F50/3D 01
3F51/FF 10
3F52/3D 01
3F53/00 00
3F54/5E 03
3F55/2A
3F56/51 ?
3F55/2A R
G
?
0300/10
0301/20
0302/30
0303/40
0304/50
0305/60
0306/70
0307/80
0308/90 R
3DEA/21 R

P 0000
0300 3E 04 05 06 07 36

L 1F2E 3DBF

?
3DD4/21 R
D/M?

```

Example 4. The user resets the computer, types a D, then types VU, bringing in the utilities package, then types RZ, switching the computer to the Z-80, running under the Z-80 monitor. He then examines location 100 up, and loads it with the sequence 10 through 90. He types a legal character to exit the load sequence, types in address 300, and examines the contents of 300 upward. He again types in a legal character, causing it to exit that mode. He then loads location 3F4F upward with the start of move, end of move, and start of destination of the Z-80 move routine, which is a part of the Challenger III utilities package. He types in a legal character to exit this routine. He then sets up a vector, or jump, to the move routine at 3F55, and types RG. The move is nearly instantaneous and returns command to the Z-80 monitor. He examines location 300 up and verifies that the move did occur. Finally the user sets up location 3DEA and goes to this routine, which switches the machine from a Z-80 to a 6800. He then verifies that he is in the 6800 and sets up the G0 vector to the Z-80, yielding the Z-80 prompter. Finally he sets up location 3DD4, which returns command to the 6502, thus showing how the user can move through to all processors.